Statistical Mechanics of Networks

TROISIEME CYCLE DE PHYSIQUE EN LA SUISSE ROMANDE

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COSIN COevolution and Self-organisation In dynamical Networks



FET Open scheme RTD Shared Cost Contract IST-2001-33555

http://www.cosin.org



• Nodes

- Period of Activity:
- Budget:
- Persons financed:
- Human resources:

6 in 5 countries April 2002-April 2005 1.256 M€ 8-10 researchers 371.5 Persons/months

EU countries Non EU countries EU COSIN participant Non EU COSIN participant

•2 Boring stuff (1/3)

- •The graph <u>size</u> is the number of its vertices.
- •The graph <u>measure</u> is the number of its edges.
- •The <u>degree</u> of a vertex in a graph is the number of edges that connects it to other vertices.
- •In the case of an oriented graph the degree can be distinguished in <u>in-degree</u> and <u>out-degree</u>.
- •Whenever all the vertices share the same degree the graph is called <u>regular</u>.
- •A series of consecutive edges forms a <u>path</u>.
 - oThe number of edges in a path is called the <u>length</u> of the path.
 - oA <u>Hamiltonian path</u> is a path that passes once through all the vertices (not necessarily through all the edges) in the graph.
 - oA <u>Hamiltonian cycle</u> is a Hamiltonian path which begins and ends in the same vertex.
 - o<u>An Eulerian path</u> is a path that passes once through all the edges (not necessarily once through all the vertices) in the graph.
 - o<u>An Eulerian cycle</u> is an Eulerian path which begins and ends in the same edge.

•2 Boring stuff (2/3)

- •Whenever all the vertices share the same degree the graph is called <u>regular</u>.
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o<u>An Eulerian path</u> is a path that passes once through all the edges (not necessarily once through all the vertices) in the graph.

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edge.





•2 Boring stuff (3/3)

- A graph is **<u>connected</u>** if a path exists for any couple of vertices in the graph.
- A graph with no cycles is a **forest**. A **tree** is a connected forest.
- The distance between two vertices is the shortest number of edges one needs to • travel to get from one vertex to the other.
- Therefore the **<u>neighbours</u>** of a vertex are all the vertices which are connected to that • vertex by a single edge.
- A <u>dominating set</u> for a graph is a set of vertices whose neighbours, along with themselves, constitute all the vertices in the graph. •
- A graph with size *n* cannot have a measure larger than mmax = n(n-1)/2. When all • these possible edges are present the graph is **complete** and it is indicated with the symbol Kn.
- The opposite case happens when there are no edges at all. The measure is 0 and the • graph is then **<u>empty</u>** and it is indicated by the symbol *En*.
- The **diameter** D of a graph is the longest distance you can find between two vertices in the graph.
- A complete **<u>bipartite clique</u>** *Ki*,*j* is a graph where every one of *i* nodes has an edge • directed to each of the *j* nodes.
- The <u>clustering coefficient</u> C is a rougher characterization of clustering with respect • to the clique distribution. C is given by the average fraction of pair of neighbours of a node that are also neighbours each other. For an empty graph En C=0 everywhere. For a complete graph Kn, C=1 everywhere.
- A **<u>bipartite core</u>** $C_{i,j}$ is a graph on i+j nodes that contains at least one $K_{i,j}$ as a • subgraph.





Router connections at small level produce a **complex** Internet structure.

•<u>2A Internet</u>

Previous maps have been computed through extensive collection of traceroutes

gcalda@pil.phys.uniroma1.it> traceroute www.louvre.fr

141.108.1.115 Rome pcpil 141,108,5,4 Unknown 193.206.131.13 Unknown rc-infnrmi.rm.garr.net 193.206.134.161 Unknown rt-rc-1.rm.garr.net 193.206.134.17 Unknown mi-rm-1.garr.net 2.1.196.25 South Cambridgesh garr.it.ten-155.net .192.37 South Cambridgesh ch-it.ch.ten-155.net 2.1.194.14 Genève geneva5.ch.eqip.net 195.206.65.105 Genève geneval.ch.eqip.net Unknown No Response 100.0.0.0 11 193.251.150.30 Unknown p6.genar2.geneva.opentransit.net 12 193.251.154.97 PARIS, FR p43.bagbb1.paris.opentransit.net



•2A Internet



Results are that we can quantify the hierarchical nature of the AS connections $P(A) \propto A^{-2}$

Plot of the C(A) show the same optimisation of the Food webs $C(A) \propto A$





skitter is a tool for actively probing the Internet in order to analyze topology and performance.



•Measure Forward IP Paths

skitter records each hop from a source to many destinations. by incrementing the "time to live" (TTL) of each IP packet header and recording replies from each router (or hop) leading to the destination host.

•Measure Round Trip Time

skitter collects round trip time (RTT) along with path (hop) data. skitter uses ICMP echo requests as probes to a list of IP destinations.

•Track Persistent Routing Changes

skitter data can provide indications of low-frequency persistent routing changes. Correlations between RTT and time of day may reveal a change in either forward or reverse path routing.

•Visualize Network Connectivity

By probing the paths to many destinations IP addresses spread throughout the IPv4 address space, skitter data can be used to visualize the directed graph from a source to much of the Internet.

http://www.caida.org/tools/measurements/skitter





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This happens at both domain and router server

•P(k) = probability that a node has k links



Faloutsos et al. (1999)



•<u>2A Internet</u>



Vazquez Pastor-Satorras and Vespignani PRE 65 066130 (2002) TABLE I. Total number of new (N_{new}) and deleted (N_{del}) nodes in the years 1997, 1998, and 1999. We also report the number of deleted nodes with connectivity k > 10.

Year	1997	1998	1999
$N_{\rm new}$	309	1990	3410
N_{del}	129	887	1713
$N_{\rm del}(k \ge 10)$	0	14	68

TABLE II. Average properties of the Internet for three different years. N, number of nodes; E, number of connections; $\langle k \rangle$, average connectivity; $\langle c \rangle$, average clustering coefficient; $\langle \ell \rangle$, average path length; $\langle b \rangle$, average betweenness. Figures in parentheses indicate the statistical uncertainty from averaging the values of the corresponding months in each year.

Year	1997	1998	1999
N	3112	3834	5287
Ε	5450	6990	10100
$\langle k \rangle$	3.5(1)	3.6(1)	3.8(1)
$\langle c \rangle$	0.18(3)	0.21(3)	0.24(3)
$\langle \ell \rangle$	3.8(1)	3.8(1)	3.7(1)
$\langle b \rangle / N$	2.4(1)	2.3(1)	2.2(1)

•2A Internet







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Why are we interested in the WebGraph?

From link analysis:

- Data mining (ex: PageRank)
- Sociology of content creation
- Detection of communities

With a "good" WebGraph model:

- Prove formal properties of algorithms
- Detect peculiar region of the WebGraph
- Predict evolution of new phenomena

Models for the WebGraph:

- Random Graph (Erdös, Renyi)
- Evolving networks (Albert, Barabasi, Jeong)
- "Copying" models (Kumar, Raghavan,...)
- ACL for massive graph (Aiello, Chung, Lu)
- Small World (Watts, Strogats)
- Fitness (Caldarelli, Capocci, De Los Rios, Munoz)
- Multi-Layer (Caldarelli, De Los Rios, Laura, Leonardi)



Albert Barabasi *Emergence of scaling in random networks* Kumar et al., *Stochastic models for the WebGraph* Broder et al., *Graph structure in the web*



•Bow-tie structure

•Small World for the SCC and the weakly connected components

Broder et al., Graph structure in the web

Cyber Communities

- **Explicit** (or "self-aware") communities:
- 1. Webrings
- 2. Newsgroup users
- 3. Gnutella, Morpheus, etc.. users
- **Implicit** communities:
- 1. Fan-Center Bipartite Cores





Kumar et al., Crawling the Web for Emerging Cyber Communities

Fractal properties

- **TUC** Thematically Unified Cluster, for example:
- 1. By content
- 2. By location
- 3. By geographical location

...and...

- 4. Random collection of websites
- 5. Hostgraph

Dill et al., Self-similarity in the web

•2C Economics and Finance

Probably the most complex system is human behaviour!

Even by considering only the trading between individuals, situation seem to be incredibly complicated.





Econophysics tries to understand the basic "active ingredients" at the basis of some peculiar behaviours. For example price statistical properties can be described through a simple model of agents trading the same stock.

"A Prototype Model of Stock Exchange" *Europhysics Letters*, **40** 479 (1997), G. C., M. Marsili, Y.-C. Zhang. Troisieme Cycle Suisse Romande Stat. Mech. of Networks-

•2C Economics and Finance

Some of the phenomena in finance can be described by means of graphs

• Stock price correlations

- •J.-P. Onnela, A. Chackraborti, K. Kaski, J. Kertész, A. Kanto http://xxx.lanl.gov/abs/cond-mat/0303579 and http://xxx.lanl.gov/abs/cond-mat/0302546
- •G. Bonanno, G. Caldarelli, F. Lillo and R. N. Mantegna http://xxx.lanl.gov/abs/cond-mat/0211546

•Portfolio composition

•D. Garlaschelli, S. Battiston, M. Castri, V. D. P. Servedio, G. Caldarelli http://xxx.lanl.gov/abs/cond-mat/0310503

•Board of Directors

•M. E. J. Newman, S. H. Strogatz and D. J. Watts, *Phys. Rev. E* 64, 026118 (2001).
•S. Battiston, E. Bonabeau and G. Weisbuch

http://xxx.lanl.gov/abs/cond-mat/0209590 (2002).

Through this new description we can

•Discover new features

•Validate Models

$$r_{i}(\tau) = \ln P_{i}(\tau) - \ln P_{i}(\tau - 1)$$

$$\rho_{i,j} = \frac{\langle r_{i}r_{j} \rangle - \langle r_{i} \rangle \langle r_{j} \rangle}{\sqrt{\left(\langle r_{j}^{2} \rangle - \langle r_{j} \rangle^{2}\right) \left(\langle r_{i}^{2} \rangle - \langle r_{i} \rangle^{2}\right)}}$$

Logarithmic return of stock *i*

Correlation between returns (averaged on trading days)

$$d_{i,j} = \sqrt{2(1-\rho_{i,j})}$$

Distance between stocks *i*, *j*

<u>A tree (a graph with no cycle) can be constructed by imposing that the</u> <u>sum of the (N-1) distances is the minimum one.</u>

Real Data from NYSE



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Correlation based minimal spanning trees of real data from daily stock returns of 1071 stocks for the 12-year period 1987-1998 (3030 trading days). The node colour is based on Standard Industrial Classification system. The correspondence is:

yellow for manufacturing green for transportation. communications. light blue for public administration black for retail trade orange for service industries "Topology of correlation based.." http://xxx.lanl.gov/abs/cond-mat/0211546 G. Bonanno, G. C., F. Lillo, R. Mantegna. Troisieme Cycle Suisse Romande Stat. Mech. of Networks-

2C Stock correlation

Data from Capital Asset Pricing Model

In the model it is supposed that returns follow

 $r_i(t) = \alpha_i + \beta_i r_M(t) + \varepsilon_i(t)$

 $r_i(t)$ = return of stock *i* $r_{\mathcal{M}}(t)$ = return of market (Standard & Poor's) α_i, β_i = real parameters = noise term with 0 mean \mathcal{E}_{i}



Correlation based minimal spanning trees of of an artificial market composed by of 1071 stocks according to the one factor model. The node colour is based on Standard Industrial Classification system. The correspondence is:

green for transportation, communications, light blue for public black for retail trade

administration

yellow for manufacturing

orange for service industries

<u>2C Stock correlation</u>



Without going in much detail about degree distribution or clustering of the two graphs We can conclude that:

the topology of MST for the real and an artificial market are greatly different. Real market properties are not reproduced by simple random models

•2C Portfolio Composition





Investors or Companies not traded at Borsa di Milano (Italy)

Companies traded at Borsa di Milano (Italy)

<u>•2C Portfolio Composition</u>



•2C Portfolio Composition



<u>•2C Portfolio Composition</u>



BEN Franklin Res Inc C Citygroup GS Goldman Sachs GBL Gabelli Asset Man LM Legg Mason INC NEU Neuberger Bergman STT State Street WM Washington Mutual

<u>•2C Portfolio Composition</u>



2C Portfolio Composition

It is not only the topology that matters.

In this case as in many other graphs the weight of the link is crucial



 $SI \propto \frac{\sum_{j} w_{ij}^{2}}{\left(\sum_{i} w_{ij}\right)^{2}}$ For every stock *i* you compute this quantity. The sum runs over the different holders • If there is one dominating holder SI tends to one • If all the holders have a similar part SI tends to 1/N For every stock *i* you compute this quantity.



For every guy *j* you compute this quantity. $HI(j) \propto \sum_{i} \frac{w_{ij}^{2}}{\left(\sum_{i} w_{il}\right)^{2}}$ The sum at the denominator runs over the different holders of *i* Then you sum on the different stocks in the portfolio This gives a measure of the number of stocks controlled

<u>2C Stock correlation</u>





•2D Food Webs

FOOD CHAIN = sequence of *predation relations* among different living species sharing the same physical space (Elton, 1927):



- -> The species ultimately feed on the abiotic environment (light, water, chemicals);
- At each predation, almost 10% of the resources are transferred from prey to predator.



A series of different interconnected food chains form a food web



•2D Food Webs

Trophic Species:

Set of species sharing the same set of preys and the same set of predators (*food web* \rightarrow *aggregated food web*).

Trophic Level of a species:

Minimum number of predations separating it from the environment.



Basal Species: Species with no prey (B) *Top Species*: Species with no predators (T) *Intermediate Species*:

Species with both prey and predators (I)

Prey/Predator Ratio = $\frac{B+I}{I+T}$

•2D Food Webs

* See Neo Martinez Group at http://userwww.sfsu.edu/~webhead/lrl.html





Pamlico Estuary (North Carolina): 14 species

Aggregated Food Web of Little Rock Lake (Wisconsin)*: 182 species → 93 trophic species

How to characterize the topology of Food Webs?





Unaggregated versions of real webs:

R.V. Solé, J.M. Montoya Proc. Royal Society Series B 268 2039 (2001)

J.M. Montoya, R.V. Solé, Journal of Theor. Biology 214 405 (2002)

Aggregated versions of real webs:



Same qualitative behaviour of their unaggregated counterparts.

We look for other quantities!.

•2D Food Webs: Spanning Trees of a Directed Graph



A *spanning tree* of a connected directed graph is any of its connected directed subtrees with the same number of vertices.



In general, the same graph can have more spanning trees with different topologies.

•2D Food Webs Spanning Trees from data

St.Martin's Island (Antilles):

44 species \rightarrow 42 trophic species 224 links \rightarrow 211 trophic links (low taxonomic resolution)

Ythan Estuary (Scotland):

134 species → 123 trophic species 597 links → 576 trophic links (taxonomic resolution : 88%)

Silwood Park (United Kingdom):

154 species → 83 trophic species 365 links → 215 trophic links (taxonomic resolution : 100%)

Little Rock Lake (Wisconsin):

182 species → 93 trophic species 2494 links → 1046 trophic links (taxonomic resolution : 31%)

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Spanning Tree:

All edges directed from level l_1 to levels $l_2 \le l_1$ are removed

•2D Protein Interactions



Network of Interaction for the protein of Baker's Yeast (Saccharomyces Cerevisiae)

How do growth and preferential attachment apply to protein networks?

 Growth: genes (that encode proteins) can be, sometimes, duplicated; mutations change some of the interactions with respect to the parent protein

 Preferential attachment: the probability that a protein acquires a new connection is related to the probability that one of its neighbors is duplicated; proportional to its connectivity

A. Vazquez et al., ComPlexUs 1, 38-44 (2003)

•2D Two-hybrid method







The two hybrid method way of detecting protein interactions

<u>•2D More refined Models</u>

With the solvation free energies taken from an exponential probability distribution $p(f) = e^{-f}$, we obtain

 $P(k) \sim k^{-2}$



The real network is random

- The detection method seesonly pairs with large enoughbinding constants
- The binding constant is relatedto the solubilities of the twoproteins
- Solubilities are givenaccording to some distribution

•2D Protein Interactions



Scale-Free Betweenness $b(k) \rightarrow$

← Scale-Free Degree distribution

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1e+06

•—• DIP •—• exponential

•2D Protein Interactions



Clustering per degree $c(k) \rightarrow$

← neighbors degree per degree Knn(k)



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